# Asymptotically Minimax Stochastic Search Strategies in the Plane

by

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# ABSTRACT

Stochastic search strategies are proposed for finding a possibly mobile target within a convex region of the plane. The strategies are asymptotically minimax as  $\varepsilon \to 0$  with respect to the time required to get within  $\varepsilon$  of the target.

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### 1. Introduction

"Princess and Monster" (1) is a zero-sum two-person game with two players restricted to a bounded, connected, two-dimensional region  $\Omega$ . The Monster (M) has maximum speed 1, the Princess (P) has maximum speed v < 1. Neither player obtains any information about the position of the other until the distance between the two is  $\leq \varepsilon$ ; at which time M captures P and the game ends. The payoff to P is the time elapsed before capture.

This game is a crude model for a surface ship M attempting to locate a submarine. Here the parameter  $2\varepsilon$  (the sweep width) is typically small relative to the dimensions of  $\Omega$ .

The P and M game is too complex to admit simple minimax strategies. Even if the continuum  $\Omega$  is replaced by a finite set of points, and even if P's strategy is known to M, M's optimal strategy can only be determined approximately, by a dynamic programming algorithm (2). Nevertheless, for convex  $\Omega$  Gal (3,4) and Fitzgerald (5) have exhibited strategies for both players that are asymptotically minimax as  $\varepsilon \to 0$ , in the sense that the ratio of the expected payoff to the minimax value approaches 1 uniformly over the opponents' strategies. They have also shown that the minimax value  $V(\varepsilon)$  satisfies

$$\lim_{arepsilon o 0}\;2arepsilon\;V(arepsilon)=|\Omega|,$$

where  $|\Omega|$  denotes the area of  $\Omega$ .

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P's strategy is easily described. Let  $Q_1, Q_2, \ldots$  be an i.i.d. sequence of random points uniformly distributed in  $\Omega$ . P starts at  $A_1$ , stays there T time units, moves to  $Q_2$  at full speed, stays there T time units, and so on. The parameter  $T \to \infty$  as  $\varepsilon \to 0$ , but  $\varepsilon T \to 0$  (e.g.,  $T = \varepsilon^{-1/2}$ ). It is not difficult to show that, no matter what strategy M uses, the expected time to capture is at least  $|\Omega|/2\varepsilon$  (approximately) when  $\varepsilon$  is small.

M's strategy is more complicated. The region  $\Omega$  is partitioned into long, narrow (width  $\varepsilon^{1/2}$ ) rectangles. M searches in one of these rectangles for a long time T (e.g.,  $T = \varepsilon^{-1/3}$ ), then moves to another and searches in it for a time T, and so on (cf. Fitzgerald (4) or Gal (5) for details).

Despite its asymptotically minimax character, this strategy for M has a defect: when  $\varepsilon$  is small, M is confined to small subregions of  $\Omega$  for very long periods of time. If the rules of the game were changed to allow P a small amount of partial information, e.g., if P were informed of the monster's position about once every  $\varepsilon^{-1/3}$  time units, then she could elude it indefinitely. Thus, the Gal-Fitzgerald strategy for M is not robust to changes in the rules which might be relevant in naval operations.

In sections 2 and 3 below, we describe alternative strategies for M that are independent of  $\varepsilon$ , asymptotically minimax as  $\varepsilon \to 0$ , and are robust to changes in the rules allowing P occasional partial information.

The proofs of the main results will be published elsewhere. They depend on new methods for studying first passages to time-dependent boundaries by certain semi-Markov processes. The results of section 2 generalize results for search in a circle obtained by Lalley and Robbins (6).

#### 2. Search in a Convex Region

Intuition suggests that a good strategy for M should produce random trajectories that are uniformly distributed in  $\Omega$ , for if this were not the case, P could gain an advantage by

hiding in the parts of  $\Omega$  less frequently visited by M.

Let  $\Omega$  be a bounded, convex region in  $\mathbb{R}^2$  with smooth boundary  $\partial\Omega$ , and let  $\nu$  be the normalized arc-length measure on  $\partial\Omega$ ,  $\int_{\partial\Omega}\ d\nu=1$ . Let  $\Theta_1,\Theta_2,\ldots$  be i.i.d. random variables such that

$$P(\Theta_i \in d\theta) = \frac{1}{2} \sin \theta \ d\theta, \qquad 0 \leq \theta \leq \pi.$$

Define a sequence of (random) points  $P_0, P_1, P_2, \ldots$  on  $\partial \Omega$  as follows: Let  $P_0$  have distribution  $\nu$ . Having defined  $P_i$ , draw the chord in  $\Omega$  from  $P_i$  that makes an angle  $\Theta_{i+1}$  with the tangent to  $\partial \Omega$  at  $P_i$  and define  $P_{i+1}$  to be the second point of intersection of the chord with  $\partial \Omega$ .

**Proposition 1.** The stochastic process  $P_0, P_1, P_2, \ldots$  is a stationary, Harris-recurrent Markov chain on  $\partial\Omega$  with stationary distribution  $\nu$ .

See Revuz (7) for the definition of Harris-recurrence.

The trajectory of M is obtained by following the chords  $P_0P_1, P_1P_2, \ldots$  in succession at unit speed. Let X(t) denote the position of M at time  $t \geq 0$ .

**Proposition 2.** The stochastic process  $X(t), t \geq 0$ , is an ergodic semi-Markov process on  $\Omega$  whose stationary distribution is the uniform distribution on  $\Omega$ . In particular, if  $f \colon \Omega \to \mathbb{R}$  is any continuous function, then

$$\lim_{t o \infty} t^{-1} \int_0^t f(X(s)) \ ds = \int_\Omega f(x) \ dx/|\Omega| \quad ext{a.s.}$$

and

$$\lim_{t o\infty} E(f(X(t))|X(0),\;\Theta_1) = \int_\Omega f(x)\;dx/|\Omega| \;\;\; ext{a.s.}$$

Let  $\Omega^0$  denote the interior of  $\Omega$ . For any  $Q \in \Omega$  define

$$au_{arepsilon} = au_{arepsilon}(Q) = \inf\{t \geq 0 \colon \operatorname{dist}\ (X(t),Q) \leq arepsilon\}.$$

Proposition 3. As  $\varepsilon \to 0$ ,

$$2\varepsilon |\Omega|^{-1}E \ au_{\varepsilon}(Q) \longrightarrow 1$$

and

$$2\varepsilon |\Omega|^{-1} \tau_{\varepsilon}(Q) \xrightarrow{\mathcal{D}}$$
 exponential with mean 1,

uniformly for Q in any compact subset of  $\Omega^0$ .

If P's strategy were to stay at a randomly chosen point of  $\Omega$ , and if this were known to M, then M could do considerably better. By following an  $\varepsilon$ -dense path through  $\Omega$  of approximate length  $|\Omega|/2\varepsilon$ , M could assure capture by time  $|\Omega|/2\varepsilon$  (approximately) and the expected capture time would be (approximately)  $|\Omega|/4\varepsilon$ . Thus our plan is only 50% efficient for locating an immobile target.

For  $\delta > 0$  let  $\mathcal{F}_{\delta}$  denote the set of continuous, piecewise continuously differentiable functions  $y(t), t \geq 0$ , valued in  $\Omega$  and such that  $|y'(t)| \leq \nu < 1$  at all t where the derivative exists and dist  $(y(t), \partial\Omega) \geq \delta$  for all t. For  $y \in \mathcal{F}_{\delta}$  define

$$\tau_{\varepsilon}(y) = \inf\{t \geq 0: \text{ dist } (X(t), y(t)) \leq \varepsilon\}.$$

**Proposition 4.** For any  $\delta > 0$ , as  $\varepsilon \to 0$ 

$$\sup_{y \in \mathcal{F}_{\delta}} |2\varepsilon| \Omega|^{-1} E \tau_{\varepsilon}(y) \to 1.$$

This shows that our strategy of following the random trajectory X(t) is almost asymptotically minimax. Since, according to the rules of the game, P is not required to stay at

least  $\delta$  away from  $\partial\Omega$ , M should actually follow the trajectory X(t) for a convex region containing  $\Omega$  in its interior, with some reasonable modification at  $\partial\Omega$  (there is no point in searching outside  $\Omega$ ). It is clear that one may construct an asymptotically minimax family of strategies for M by using Proposition 4.

Our strategy does not have the "localization" property of the Gal-Fitzgerald strategy. Even if P is given the position and direction of M from time to time, she will not be able to predict its course for very long, in view of Proposition 2. Therefore, the strategy for M that we have described is not only fully efficient in the minimax sense, but also robust to partial information.

# 3. Search in a Parallelogram

A somewhat different search plan may be used when  $\Omega$  is a parallelogram. Like that of the preceding section, this strategy does not suffer from the localization defect. For simplicity of exposition we shall describe the strategy for the square  $\Omega = [0,1] \times [0,1]$ .

Consider the torus  $\hat{\Omega} = \mathbb{R}^2/2\mathbb{Z}^2$ , and let  $\pi \colon \mathbb{R}^2 \to \hat{\Omega}$  be the natural projection mapping. The torus  $\hat{\Omega}$  may also be thought of as the square  $[0,2] \times [0,2]$  with opposite sides identified. Consider the mapping  $\xi \colon [0,2] \times [0,2] \longrightarrow [0,1] \times [0,1]$  defined by

$$\xi(x,y) = (x,y)$$
 if  $0 \le x \le 1$ ,  $0 \le y \le 1$ ;  
 $= (2-x,y)$  if  $1 \le x \le 2$ ,  $0 \le y \le 1$ ;  
 $= (2-x,2-y)$  if  $1 \le x \le 2$ ,  $1 \le y \le 2$ ;  
 $= (x,2-y)$  if  $0 \le x \le 1$ ,  $1 \le y \le 2$ .

Since  $\xi$  maps corresponding points on opposite sides of  $[0,2] \times [0,2]$  onto the same point of  $[0,1] \times [0,1]$ ,  $\xi$  may be projected to a mapping  $\xi$ :  $\hat{\Omega} \to \Omega$ . The composition  $\xi \circ \pi$ :  $\mathbb{R}^2 \to \Omega$  is a continuous mapping of the plane onto the unit square.

Fix  $0 < \rho < 1$ . Let  $P_1, P_2, \ldots$  be i.i.d. random vectors in  $\mathbb{R}^2$  with the uniform distribution on the circumference of the circle of radius  $\rho$  centered at (0,0). Let  $P_0$  be

uniformly distributed on  $[0,1] \times [0,1]$ , and independent of  $P_1, P_2, \ldots$  Define a random path Y(t) in  $\mathbb{R}^2$  as follows. Start at  $P_0$ , move at unit speed along the line segment from  $P_0$  to  $P_0 + P_1$ , then move at unit speed along the line segment from  $P_0 + P_1$  to  $P_0 + P_1 + P_2$ , and so on.

Our search plan calls for M to follow the projection onto  $\Omega$  of the random path Y(t), so that the position of M at time t is

$$X(t) = \xi(\pi(Y(t))).$$

**Proposition 5.** The stochastic process X(t),  $t \geq 0$ , is an ergodic semi-Markov process on  $\Omega$  whose stationary distribution is the uniform distribution on  $\Omega$ . In particular, if  $f \colon \Omega \to \mathbb{R}$  is continuous, then

$$\lim_{t o\infty} \; rac{1}{t} \; \int_0^t \; f(X(s)) \; ds = \int_\Omega \; f( ilde x) \; d ilde x \quad ext{ a.s.}$$

and

$$\lim_{t o\infty}\; E(f(X(t))|P_0,P_1) = \int_\Omega\; f(x)\; dx \quad ext{ a.s.}$$

Let  $\mathcal{F}_{\delta}$  be defined as in section 2. For  $z \in \mathcal{F}_{\delta}$ , let  $\tau_{\varepsilon}(z) = \inf\{t \geq 0 : \text{ dist } (X(t), z(t)) \leq \varepsilon\}$ .

**Proposition 6.** Let  $0 < \delta < 1/2$ . If  $\rho < 2\delta$ , then as  $\varepsilon \to 0$ ,

$$\sup_{y \in \mathcal{F}_{\delta}} \ 2\varepsilon E \ \tau_{\varepsilon}(z) \to 1.$$

## 4. Concluding Remarks

(1) We conjecture that any ergodic semi-Markov process with uniform stationary distribution may be used to obtain asymptotically minimax strategies for M.

- (2) The boundary effect implicit in Propositions 3, 4, and 6 deserves further study. How should M behave near  $\partial\Omega$ ?
- (3) It would be useful to have a good definition of a modified search game when the boundary  $\partial\Omega$  is fuzzy.
- (4) It would be useful to make a systematic study of modifications of the game in which one or both players is allowed partial information about the movement of the other.

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